# Scaling Databases to Infinity and Beyond! (Well, almost Infinity!)

NoCOUG – August 2021

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## Agenda

- 1. Intro
- 2. Tenets
- 3. Keys
- 4. Tables
- 5. Access Path
- 6. DB Engines

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## Intro

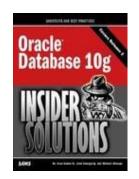


## About the Speaker

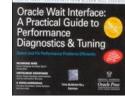
- Currently Sr. Database/Data Architect @ PayPal
- Has been working with Oracle Databases and UNIX for 3+ decades
- Working on various NoSQL/Big Data technologies for the past 6 years
- Design and Implement High scale systems Both Oracle and NoSQL
- Author, Technical editor, Oracle ACE Alumni, Frequent speaker
- Loves to mentor new speakers and authors!
- http://www.linkedin.com/in/johnkanagaraj









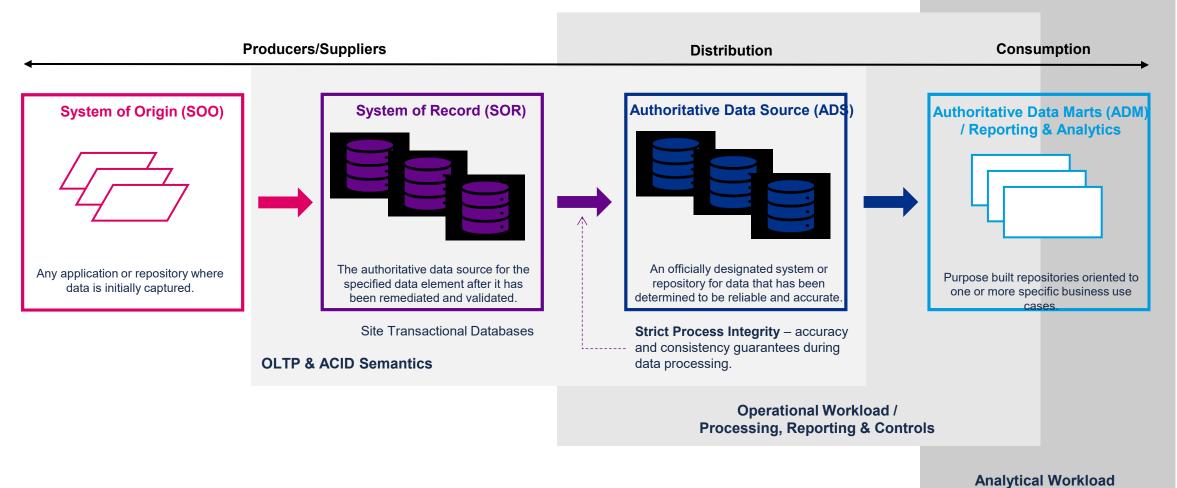




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## **Definitions: Data Platforms**



Write Once Read Many (WORM)



#### **Site Data Architecture**

#### System of Record (SOR)



The authoritative data source for the specified data element after it has been remediated and validated.

Site Transactional Databases

Authoritative Data Source (AD



An officially designated system or repository for data that has been determined to be reliable and accurate.

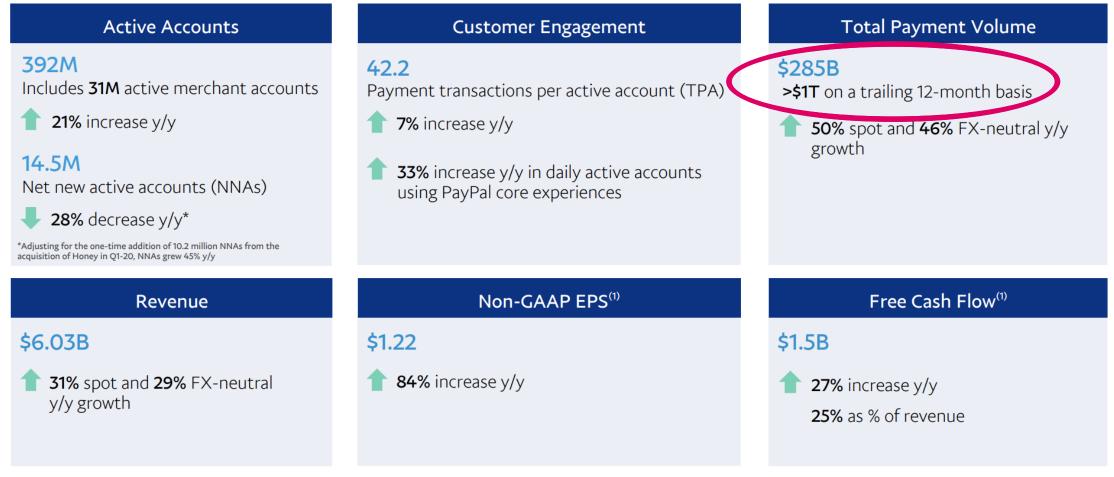
Strict Process Integrity – accura and consistency guarantees during data processing.

## Core Tenets



## Know [Data About] Your Data (KYD)

### First Quarter 2021 Summary Strong performance across key performance metrics





## Site Data Architecture

Core Tenets





## **Challenges at Scale**

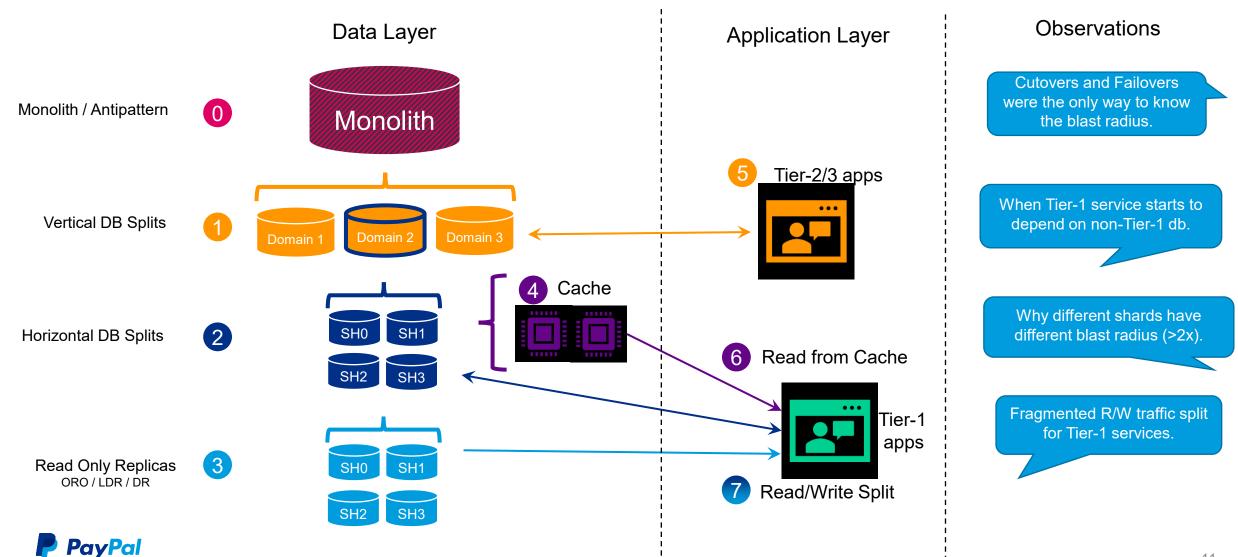
- Pushing the limits
  - Connections
  - Memory
  - Interconnect
  - CPU
  - DDL on busy tables
  - RAC reconfiguration
  - Redo rate
  - I/O latencies
  - SAN Storage limits
  - Replication latencies
  - HA requirements



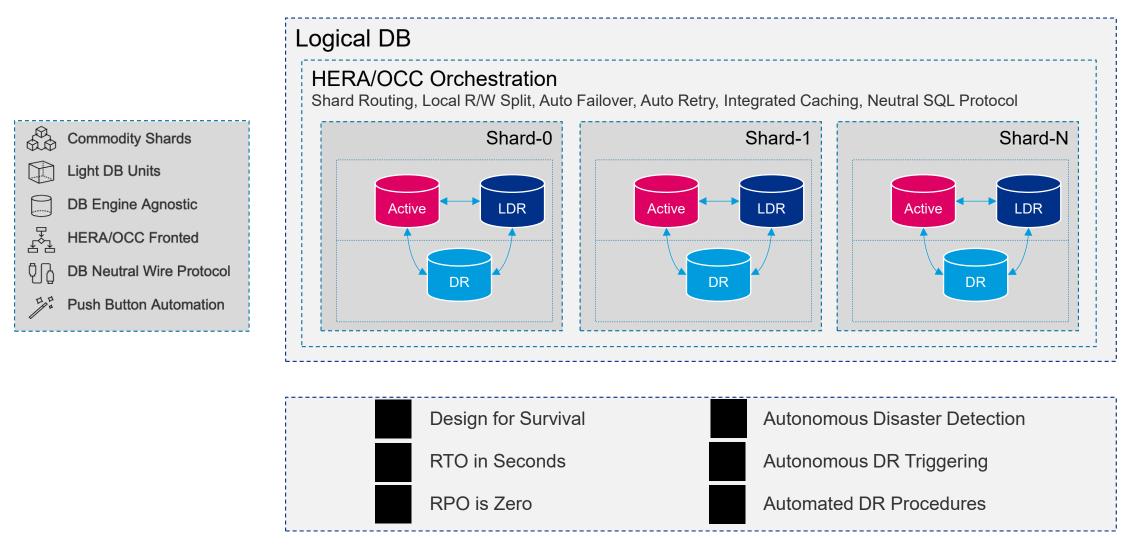
- Solutions
  - Custom Connection pooling and multiplexing (OCC)
  - Read Scale out (replication)
  - Microservice oriented architecture (logical separation)
  - Custom HA caching (Juno)
  - Custom Sharding
  - Active-Active operation using Oracle RO's and GoldenGate
  - Storage Tiering and Archiving

  - Moving to Cloud! ©

## Scaling and Blast Radius Resiliency Patterns



## **Resiliency Vision**



### PayPal

## Keys: Timed UUID



## Timed UUID / PayPal Variant

UUID Version 1- timed; 4 - random

128-bit number in hexadecimal format:

6E8BC430-9C3A-11D9-9669-0800200C9A66 [time-low]-[time-mid]-[version-and-time-high]-[clock-misc.]-[node]

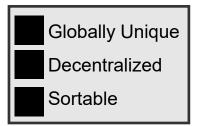
count of 100 nanosecond intervals Since 10/15/1582 00:00:00.00 UTC

 Decimal equivalent:
 [SS.FFFF]-[HH24:MI]-[YYYY-MM-DD]
 [50.123456]-[11:20]-[2021-05-27]

 Decimal equivalent unshuffled:
 [YYYY-MM-DD]-[HH24:MI]-[SS.FFFFF]
 [2021-05-27]-[11:20]-[50.123456]

Hexadecimal unshuffled:

11D9-9C3A-6E8BC	<b>430</b> -9669-
0800200C9A66 4 bits 60 time-based bits	64 random bits



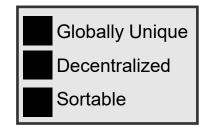
Primary Key, Idempotency Key, Time-Based Partition Key



## Timed UUID / PayPal Variant / Collision Rate

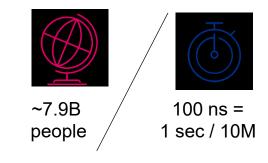
Hexadecimal unshuffled

11D9-9C3A-6E8BC4	30-9669-
080F210C9A66 4 bits 60 time-based bits	64 random bits



To allow for 50% probability of one collision,

we need to generate 9.1x10<sup>9</sup> UUIDs within 100 nanoseconds (ns).



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## KYD: Table Categories



## Table Categories

	Immutable	Mutable
	Reference or configuration data.	Standard business objects like merchant, customer, customer account, customer address, etc
Master Data	Country Lookup	Customer Balance Merchant

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## Table Categories

	Immutable	Mutable
Master Data	Reference or configuration data.	Standard business objects like merchant, customer, customer address, etcCustomerCustomer BalanceMerchant
Transactional Data	Customer Activity LogJournal	Time-based events that are modifiable until they 'close', such as payments or customer cases.PaymentCustomer Case

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## **Hybrid Tables**

### [Master + Transactional Data]



[Do Not Mix]

## KYD: Access Path



## **Problem Statement**

#### **1. Customer Activities**

Acct #	Time	Attribute 1	Attribute 2	Attribute 3	Attribute N	PK-	ACCOUNT#
100	t1	А	Х	6	!@	FN	TIMED_KEY
400	t1	А	Y	7	#\$		Attribute_1
300	t2	В	Х	8	%^		Attribute_2
100	t3	С	Z	0	^&		Attribute_3
300	t4	D	Z	9	*(		
							Attribute_N

#### 2. Query Requirements

Customers need to query and search their data.



#### CUSTOMER\_ACTIVITY

## Problem Statement cont.



#### **1. Denormalized Table**

#### CUSTOMER\_ACTIVITY (CA)

ACCOUNT#

TIMED KEV

PK

#### 2. Query Requirements

Need to search and query within ACCOUNT# across all attributes

#### 3. Required Indexes

L						
	Attribute_1		Attribute_1		IDX1	Account#, Attribute_1
	Attribute_2		Attribute_2		IDX2	Account#, Attribute_2
	Attribute_3		Attribute_3		IDX3	Account#, Attribute_3
	Attribute_4		Attribute_4		IDX4	Account#, Attribute_4
Payload _	Attribute_5		Attribute_5		IDX5	Account#, Attribute_5
/ Value	Attribute_6		Attribute_6		IDX6	Account#, Attribute_6
		-		·		
			Attribute_N		IDX N	Account#, Attribute N
				l	—	, <u> </u>
	Attribute_N					

## **Denormalized Index**

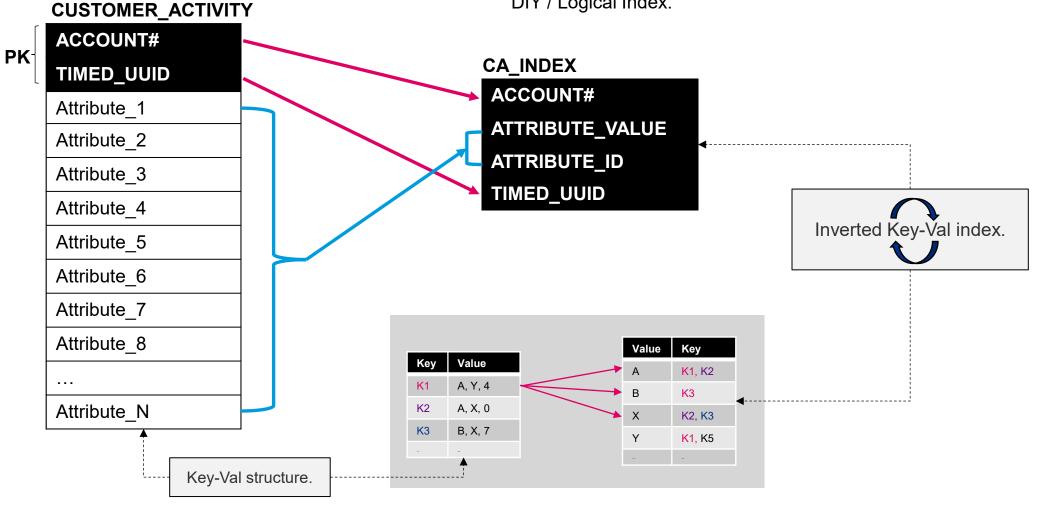
**1. Denormalized Table** 

(Denormalized)<sup>2</sup>



#### 2. "Denormalized" Index

DIY / Logical Index.

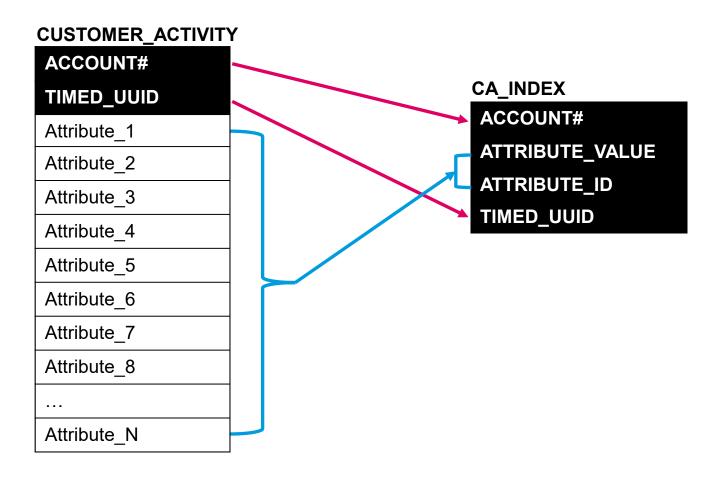


## Query



#### 1. Denormalized Table

2. "Denormalized" Index



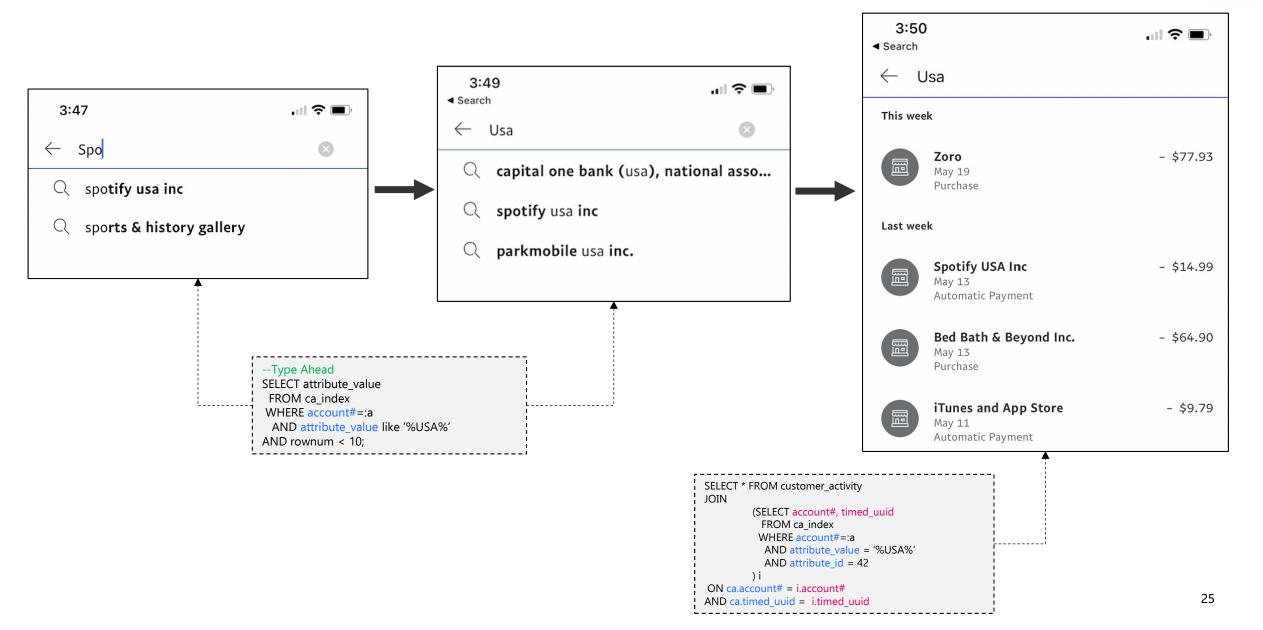
#### 3. How does it work?

Phase 1: Query Index Table Phase 2: Fetch from CA

[	
	-pseudo code
1 7	SELECT *
¦ F	ROM customer_activity
¦ J	OIN
!	(SELECT account#, timed uuid FROM
	ca index
!	WHERE account#=:a
	AND attribute_value = 'Arch'
	—
	AND attribute_id = 42
	) I
	ON ca.account# = i.account#
A	AND ca.timed uuid = i.timed uuid
	-dual PK access
Ĺ	

### Search



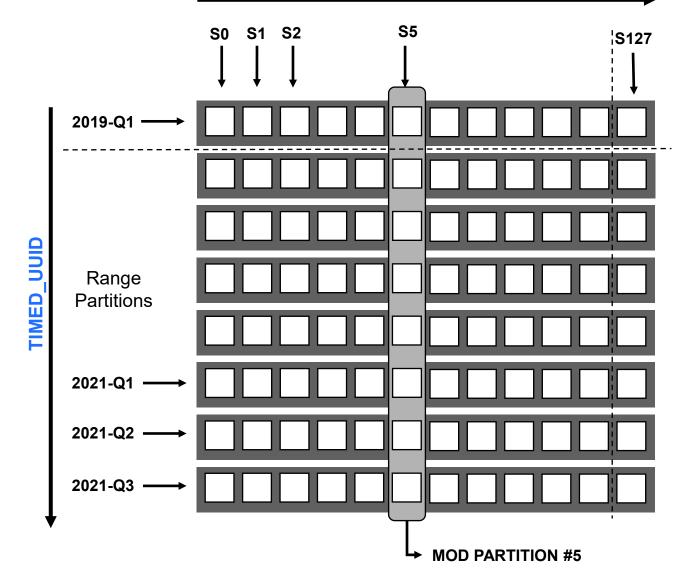


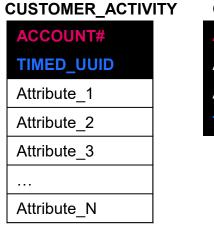
## Scaling it Out

Sharding / Scale-Out











ACCOUNT# = Shard Key TIMED\_UUID = Range Key

PK on (ACCOUNT#, TIMED\_UUID)

#### Current State (7+ y/o CAM)

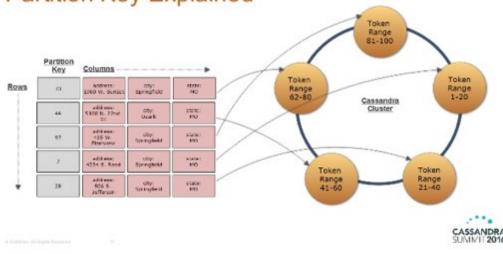
- 8 Physical Shards
- ~100TB/Shard
- >billions of reads/writes/day/table

## **DB Engine Agnostic**

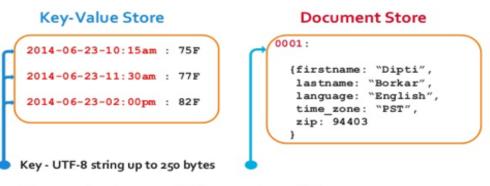
#### Why use NoSQL?

#### Flexibility

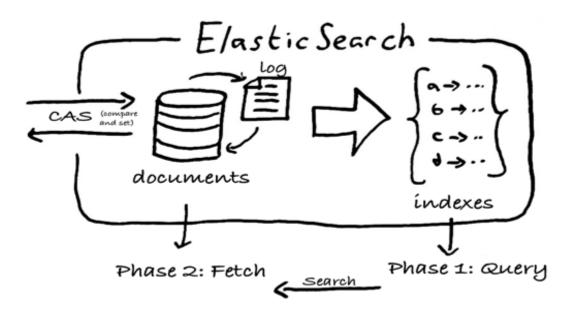
- Flexible schemas that enable faster and iterative development.
- Ideal for semi-structured and unstructured datasets. **Scalability**
- Designed to scale out by using distributed clusters of hardware.
- Some cloud providers manage it behind-the-scenes **Performance**
- Optimized for specific data models (e.g., document, keyvalue)
- Optimized for access patterns.
   Partition Key Explained



#### Couchbase can act as a



Value - can be o bytes – 20 MB (best practice < 1 MB)





# Appendix



## Query vs Search

	Query Engine	Search Engine (ES)
1.Input	Know exactly what you are looking for.	Exact value is not required for searching. Supports fuzzy, partial, proximity, etc., match.
2. Output	Returns only results that match.	Top-N ranked matches based on <u>relevance</u> scoring using <u>tf-idf</u> .
	Completeness and accuracy is guaranteed.	Do not need to retrieve all results. First few pages, OK.
3. Data Access	Predominately index-based access. Primary or secondary index based.	Query each shard in the index (Phase 1). Populate local priority queue (top-n results). Combine into global queue (global top-n). <u>Fetch</u> top-n documents (Phase 2).
4. Performance	Optimal performance / access path.	Search has much more work to do.
5. Service Time	Supports millions of executions per second. From sub-msec.	Takes longer than regular query. 10-20 msec - considered good. 100-200 msec - under heavy load. "Depending on the <u>search complexity</u> (term vs phrase vs proximity), it can be 10 to 20 times longer than simple term search."
6. Core Strengths	Predictable / Systematic Lookups.	Interactive (Human!) Search / Complex Investigation.

### PayPal

## Key-Val and Inverted Indexes



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### **Composite Key-Val Structures**

	Data Ta	able
Acc#	Timed Key	Value
ABC	TK1	A1, B3, C4
ABC	TK2	A1, C2, X0
ABC	ТК3	C2, X0, Y4

#### Index Table

