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SQL Performance Hero and OMG Method Play the Anti-patterns Greatest Hits

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Survey Says

- DBAs
- Developers
- Architects
- Heavily non-Oracle development shop
- Concerned with performance
- Access to production size database
- Easy access to running traces, Enterprise Manager



Introduction to OMG Method

- OMG Method focuses on
 - Refactoring SQL
 - Indexing
 - Refactoring application side code
- OMG Method targets performance problems created by Developers Inexperienced in Oracle technologies (DIO)
- OMG Method requires (almost) no DBA privileges other than indexing
 - No tracing, (almost) no init.ora parameter changes

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Fair Warning

- No demos
- No "proofs"
- Quick fixes

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Requirements for SQL Performance Heroes

- Good SQL fundamentals
- Able to read basic explain plans
- Understand basic performance statistics from autotrace
- Courage to make and test changes
 - Don't take my word for it!
- Willingness to work with and educate DIOs

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Why OMG Method

- Vast majority of performance problems are result of DIOs'
 - Lack of training in SQL and Oracle
 - Lack of interest in SQL and Oracle
 - Misinformation about SQL and Oracle performance
 - Resistance to PL/SQL
 - Focus on OO, procedural and functional programming techniques
 - Iterative thinking vs set thinking

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Anti-Patterns

- Definition
 - Common SQL or design practice that results in poor performance
- OMG Method identifies common anti-patterns and techniques to fix them
 - Always verify that OMG fixes actually improve performance
- OMG Method does not address schema design problems
 - No changes to tables or columns

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Understanding Common Design and DIOs Anti-patterns

- Overly Generic Data Models
 - OBJECT, INSTANCE, ATTRIBUTE, ATTRIBUTE_VALUE structures
- Fat, Unnormalized Tables
 - Often with in-line CLOBs
- Fear of Joins
 - "Joins are to be avoided a all costs" mentality
- Failure to Understand SQL query cost in application code
- Iterative vs Set World View

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Understanding Common Design and DIOs Anti-patterns

- Unmanaged Surrogate Primary Keys
 - (Nearly) all tables have surrogate primary keys
 - Values for same row is not consistent across enviroments, e.g., COMPANY_ID value for same company differs across production, development, test environments
 - Typically use additional *lookup* columns
- Widespread use of *Dummy* values instead of NULL
 - DIOs uncomfortable using NULL
 - Misunderstanding of performance issues with NULL

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Understanding Common Design and DIOs Anti-patterns

- "Indexed searches are always better"
- Lack of documentation

Avoid Dynamic SQL

- Avoid/eliminate dynamic SQL, e.g. creation and execution of SQL queries created by concatenating strings
 - Particularly problematic when using literals for constants
- Use prepared statements with bind variable
- Dynamic SQL results in heavy parsing overhead and SGA memory usage
 - Child cursors may be created even if the only differences between SQL queries is literal values
 - Potential for SQL Injection



Inline Views

• In SQL code, an inline view is a subquery used in place of a table, e.g.,

```
SELECT ...
FROM
(SELECT ...)
```

Avoid/Replace Materialized Inline Views

- Inline views typically results in an "inline view" being created in the execution plan
 - Referred to as materialized inline view (MIV)
- Oracle may also merge the SQL inline view with the outer query
- MIVs produce a *result set,* e.g., a temporary table (not to be confused with *Global Temporary Table*)
 - MIVs are never indexed
 - Joins with a MIV effectively perform a Full Table Scan (FTS) against the MIV, e.g. *multiple FTS!*
 - Poor performance if result set is large



Avoid/Replace Materialized Inline Views

- DIOs frequently write inline views which can and should be replaced by joins
 - Generally can be done with little or no understanding of underlying schema semantics
 - Try /*+ MERGE */ hint first; generally doesn't improve performance, but worth trying
 - May also help in rewrite

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Merged Inline Views

- As the Cost Based Optimizer has evolved, it frequently merges SQL inline views with the outer query
- Frequently not a performance improvement!
 - Particularly with poorly written SQL inline views
 - 10G's merging is much better than 9i's
- Try /*+ NO_MERGE */ hint

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Never Update Primary Key Columns

- Primary key (PK) columns should never be updated, even to current value
- Common DIO approach is to update all columns in a row
- Updating PK columns forces examination of referencing foreign key (FK) constraints on child tables
 - General performance issue, even if FK columns indexed
 - Results in FTS if FK columns not indexed

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Avoid/Remove Unnecessary Outer Joins

- DIOs frequently add outer joins "just to be safe"
- Outer joins may be expensive, limiting CBO choices
 Be sure join columns are indexed
- Work with developer or end user to determine if outer join is needed

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EXISTS vs IN

- Replacing IN with EXISTS often produces dramatic performance improvement
- IN uses uncorrelated subquery
- SELECT ...

```
FROM table_1 outer
WHERE
outer.col_1 IN
 (SELECT inner.col_1
 FROM table_2 inner
 [WHERE ...])
```

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IN Performance Issues

- IN may perform poorly
 - Produces result set, effectively a materialized inline view
 - Result set is unindexed
 - Result set is scanned for every row in outer query
 - Large result set is well known performance killer
- IN should only be used when the result set is small
- Note that if the value of outer.col_1 is NULL, it will never match the result of the IN
 - Use NVL on both the inner and outer columns if NULL must be matched

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EXISTS vs IN

- DIOs seldom know how to use EXISTS as it involves a correlated subquery, e.g., a join between column(s) in the outer and column(s) in the inner query
- Replace the uncorrelated subquery with a subquery by joining the outer column from the IN clause with an appropriate column in the subquery

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EXISTS Correlated Subquery

• SELECT ...

```
FROM table_1 outer
WHERE
EXISTS
(SELECT 'T' -- use a simple constant here
FROM table_2 inner
WHERE
outer.col_1 = inner.col_1
[AND ...]) - WHERE predicates from original
query
```

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EXISTS Correlated Subquery

- The join columns (inner.col_1 in example) from the table in the correlated subquery should be indexed
 - Check to see if appropriate indexes exist; add them if needed
- Use a constant in the SELECT of the correlated subquery; do not select the value of an actual column
- Note that SELECT DISTINCT is unnecessary for both IN and EXISTS



Subquery Factoring using WITH

- Very powerful (and virtually unknown)
- Many DIO written queries use *identical* subqueries/inline views repeatedly
- Often lengthy UNIONs

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Often lengthy UNIONs

```
SELECT ...
FROM
   table_1,
   (SELECT ... FROM table_2, table_3, ... WHERE
     table_2.id = table_3.id) IV
   WHERE ...
UNION
   SELECT ...
   FROM
   Table_4,
   (SELECT ... FROM table_2, table_3, ... WHERE
     table_2.id = table_3.id) IV
   WHERE ...
   UNION ...
```



Performance Issue

- Oracle's CBO is not aware of identical nature of subqueries (unlike programming language optimizers)
 - Executes each subquery
 - Returns distinct result set for each subquery
 - Redundant, unnecessary work

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Subquery Factoring

- Subquery factoring has two wonderful features
 - Generally results in improved performance
 - Always simplifies code
 - Factored subquery only appears once in the code as a preamble
 - Referenced by name in main query body
 - More readable, easier to maintain and modify

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Syntax

/* Preamble, multiple subqueries may be defined */
WITH

```
(SELECT ...)
```

```
AS pseudo_table_name_1
```

```
[, (SELECT ...) AS pseudo_table_name_2 ...]
```

/* Main query body */

SELECT ...

```
FROM pseudo_table_name_1 ...
```

```
... -- typically UNIONs
```

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Example

```
    Applying this to the example

/* Preamble */
WITH
    (SELECT ... FROM table_2, table_3, ... WHERE table_2.id
      = table_3.id
   AS IV
/* Main query body */
SELECT ...
   FROM
   table_1, IV - IV is pseudo table name
   WHERE ...
UNION
   SELECT ...
   FROM
   Table_4, IV
   WHERE ...
UNION ...
```



Factoring Options

- Oracle may perform one of two operation on factored subqueries
 - Inline performs textual substitution into main query body
 - Effectively same query as pre-factoring
 - No performance improvement due to factoring
 - Still more readable
 - Materializing factored subquery
 - Executes the factored subquery only once
 - Creates true temporary table (not Global Temporary Table)
 - Populates temporary table with direct load INSERT from factored subquery

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Materialized Factored Subquery Issues

- Materialized Factored Subqueries (MFS) issues CREATE TABLE for temp table at least once (on 1st execution)
- Multiple tables may be created if query executions overlap
- Tables are reused if possible
- Recursive SQL performs INSERT
- Data is written to disk
- Doesn't always result in performance improvement



Hints for Subquery Factoring

- /*+ Materialize */ will force materializing
 - Seldom, if ever, needed
 - Oracle only materializes when subquery used more than once (but verify)
- /*+ Inline */ will force textual substitution
 - Use when materializing does not improve performance
- Other hints may be used in factored subquery, e.g. INDEX, etc.
 - Note that MERGE and NO_MERGE may be combined with INLINE
- Hint follows SELECT in factored subquery
 - WITH (SELECT /*+ hint */ ..) AS ...

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INDEX Hints

- DIO often believe everything should use indexes
- Frequent use of *unqualified INDEX* hint, e.g., only table name specified, but no index
 - SELECT /*+ INDEX (table_name) */
 - Yes, this does work!
- Oracle will always use an index, no matter how bad
 - Unclear which index will be used; documentation says "best cost", but unclear if true; experience suggests 1st in alphabetical order
 - Further complicated by poor indexing
- Fix by either
 - Qualifying hint by specifying index name(s)
 - Removing hint entirely
 - Removing the hint often improves performance



Constant Data Conversion Issues

- When comparing a VARCHAR2 (or CHAR) column to a constant or bind variable, be sure to use string data type or conversion function
- Oracle does not always do what you would expect
 - WHERE my_varchar2_col = 2 does not convert 2 to a string!!!
 - Converts every rows's my_varchar2_col to a number for the comparison
 - Generally results in FTS
 - Common cause of "I just can't get rid of this FTS"
- Common problem with *overloaded* generic and OO models
- Even SQL Performance Heroes get bit!!!



Eliminate Unnecessary Lookup Joins

 Tables with unmanaged surrogate keys typically have lookup/alternate key column(s) with consistent data across environments

- Very common with generic and OO models

- Typical code is:
- SELECT

```
FROM child_table, reference_table
WHERE
    child_table.reference_table_id =
reference_table.reference_table_id
    and reference_table.lookup_column =
`constant'
```

Results in access to reference_table for every applicable row in child_table



Eliminate Unnecessary Lookup Joins

- Even worse when UPPER/LOWER function applied to lookup_column (unless appropriate functional index exists)
- Replace with scalar subquery

Only performs scalar subquery once



Improving Pagination

- Pagination refers to returning row n through m from an ordered result set using ROWNUM
 - Typically for data on a web page or screen
- Common, worst case code:

```
SELECT t1.col_1,...

FROM

(SELECT *

FROM table_1

WHERE ...

ORDER BY ...) t1

WHERE

ROWNUM between n and m
```

Improvement Steps

- 1. Replace literals with bind variables
- 2. Replace "*" in innermost inline view with desired columns
 - Potentially reduces unnecessary I/O and sort processing
- Refactor the query so that inline view only returns 1st m rows and use /*+ FIRST_ROWS */ hint (per Tom Kyte's *Effective Oracle by Design* on *Pagination with ROWNUM*)

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Improvement Step #3

```
SELECT *
FROM
  (SELECT /*+ FIRST_ROWS */
    ROWNUM AS rnum, a.*,
    FROM
    (SELECT t1.col_1,...
    FROM table_1
        WHERE ...
        ORDER BY ...) a
    WHERE
        ROWNUM <= :m)
WHERE rnum > = :n
```

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Improvement Step #4

- Replace the columns in innermost inline view with ROWID and join to table in outermost query
 - May provide substantial I/O performance improvements on fat tables, particularly those with inline CLOBs

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Improvement Step #4

```
SELECT t1.col_1,...
   FROM
   table 1,
   (SELECT /*+ FIRST_ROWS */
       ROWNUM AS rnum, inner_row_id
       FROM
           (SELECT ROWID inner_row_id -- innermost query
            FROM table 1
                WHERE ...
                ORDER BY ...)
           WHERE
            ROWNUM \langle = : m \rangle
       WHERE rnum > = :n
       AND table_1.ROWID = inner_row_id
```

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UPDATE and DELETE Performance

- "I'm DELETEing/UPDATEing a few rows. It's virtually instantaneous when I test it in my development environment, but takes a very long time in production!" – Joe the DIO
- Check for indexes on FK constraint columns of child tables.
 - Lack of indexes on FK constraints requires an FTS of each child table for each row to be DELETEd/UPDATEd in parent table
 - Common problem with history tables
- Add appropriate indexes



UPDATE and DELETE Performance

- Look for foreign key constraints using Cascade Delete
 - Hierarchy of cascade deletes can result in very poor performance
 - Unclear if circular references ever complete
- Beyond scope of OMG
 - Application code may depend on existence of Cascade Delete
 - Quick fix may be temporarily altering constraints

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Eliminate *Bitmap Conversion to/from ROWIDs* Execution Plan Step

- Known performance problem unrelated to existence/use of bitmap indexes
- May be huge CPU hog
- alter session set "_b_tree_bitmap_plans" = false
- Default in 9i and 10G is true
- Real fix requires setting init.ora parameter
- _b_tree_bitmap_plans = false
- Requires database restart; can't be set dynamically

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Add Indexes on Foreign Key Constraints

- FK constraints should always be indexed
 - Have not yet seen exception to this rule (but always interested)
- Primary performance gains
 - Improved join performance fundamental feature of CBO
 - UPDATE and DELETE performance
 - Oracle apparently still performs table level locks, despite statements to contrary

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Add Foreign Key Constraints

- "FK constraints hurt performance. We'll enforce referential integrity (RI) in the application" – Flo the DIO
 - Translation: "We won't make any mistakes in the application code"
 - Won't really verify RI in the application
 - True verification would result in worse performance
- It doesn't matter how well the system performs if the data is corrupt!
 - Earned big \$ as expert witness demonstrating issues with lack of FK constraints
- CBO uses existence of FK constraints
- Adds to effective documentation of system

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Eliminate Redundant Indexes

- Redundant indexes, e.g., indexes with identical leading columns
 - Common DIO anti-pattern
- Impacts INSERT/UPDATE/DELETE performance
- Confuses CBO
 - Unclear how CBO selects index when two (or more) have needed leading columns, but different trailing columns
- Rules of thumb
 - Eliminate index with most trailing columns
 - Indexes with more than 3 columns are suspect
 - PK indexes with trailing columns should be reduced to PK only

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Reduce Unnecessary and Redundant Queries

- Worst real world case
 - 80,000 individual queries from application takes 3+ hours
 - Single query took under 30 seconds
- Individual query is not performance problem
 - Total number of queries is problem
- Two general cases
 - 1. Iteration
 - DIO issues large number of SELECTs, typically performing join, calculations or sorts in application
 - Generally easy to replace with single query
 - 2. Redundant Queries
 - DIO issues same query repeatedly for unchanging data, typically refreshing page/screen, i.e., field label
 - Requires changes to application code structure
 - Not usually Hero's domain



Add Appropriate Functional Indexes

- Functional indexes (FI) are great quick fixes for many anti-patterns
- Two most common anti-patterns

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Mixed case string columns

- Column contains mixed case data used for both lookup/filtering and display
 - Good design would be two columns, one for lookup and one for display
- (Somewhat) knowledgeable DIO uses UPPER(column_name)
 - Less knowledgeable use LOWER(column_name)
- Add appropriate index(es)
 - If possible, standardize queries to use one function
 - May need to add both indexes :-{

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Eliminating Dummy Values

- DIOs typically use dummy values in place of NULL, e.g., -99
- Queries use: WHERE column_name <> -99 instead of WHERE column_name IS NOT NULL
- <> kills use of index on column_name
- If significant percentage of rows contain dummy value, add functional index to improve performance
 - NULLIF(column_name,-99)
- Queries need to be modified to use function
 - WHERE NULLIF(column_name,-99) IS NOT NULL
- Real world cases may involve multiple dummy values, e.g. -9, -99 and -999 (really!)
 - Use DECODE, CASE or other function



Use PL/SQL for Bulk Operations

 Use of BULK COLLECT and FORALL provides huge performance improvements over application side operations

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Summary

- Many anti-patterns easily identifiable
- Many anti-patterns subject to easy, quick and safe fixes
 OMG Tips won't work for every query
- SQL Hero needs to be willing to modify queries and test results
- SQL Hero needs to understand why DIOs use antipatterns and educate them

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Q&A





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