

# Don't Be In a Funk: Use Analytic Functions



Philip Rice  
University of California Santa Cruz



# Overview

- My involvement came from a performance problem
- Overlap with some of the “traditional” aggregate functions: e.g. max, avg, count  
-- same keyword, similar syntax
- Goal today: raise awareness of possibilities, know when to consider as an option; not a comprehensive view of all functions



# Overview

- Functions are used in SELECT statement:  
likely to be most helpful in reporting situations
- Better functionality: traditional approach can be much more difficult or nearly impossible
- Performance improvement is likely to be more obvious with larger datasets, difference can be hours down to minutes(!)



# General Syntax

- Function(arg1, ..., argn) **OVER** (  
[PARTITION BY <...>] [ORDER BY <...>]  
[window\_clause] )
- “**OVER**” is indicator of analytic function
- PARTITION BY is comparable to GROUP BY
- window\_clause is not as commonly used, but can be helpful, e.g. looking at different time periods on same row of output (examples later)
- window\_clause (partial) syntax is  
[ROW or RANGE] BETWEEN <start> AND <end>



# Example: “traditional” count

- ```
select count(*), OBJECT_TYPE
from all_objects
where owner = 'OUTLN'
group by OBJECT_TYPE;
```

```
COUNT(*) OBJECT_TYPE
-----
4 INDEX
1 PROCEDURE
3 TABLE
```

- Non-aggregated columns must be in GROUP BY
- What if we want to show detail at same time as the aggregate?



# If “OVER” is empty, acts on whole set

- ```
select object_name,object_type,
       count(*) OVER () tot_count,
       count(*) OVER (PARTITION BY object_type) type_count
from all_objects where owner = 'OUTLN';
```

OBJECT_NAME	OBJECT_TYPE	TOT_COUNT	TYPE_COUNT
OL\$NAME	INDEX	Total	4
OL\$HNT_NUM	INDEX	for ALL	Total
OL\$SIGNATURE	INDEX	rows,	for
OL\$NODE_OL_NAME	INDEX	on	each
ORA\$GRANT_SYS_SELECT	PROCEDURE	each	object
OL\$NODES	TABLE	detail	type
OL\$HINTS	TABLE	line	
OL\$	TABLE		



# Timing of execution in SQL

- Analytic functions are computed after:
  - All joins
  - WHERE clause
  - GROUP BY
  - HAVING
- Main ORDER BY of query is *after* analytic function
- So AFs can only appear in select list and in main ORDER BY clause of query



# Timing of execution in SQL

## ➤ Stages:

Joins, WHERE, GROUP BY, HAVING



Apply analytic functions to result set rows (aka “partition”)



Apply ORDER BY clause  
(from main query)





# Prep for ROW\_NUMBER and RANK

- Select object\_type “ObjTyp” to\_char(last\_ddl\_time, 'yyyymmdd hh24miss') last\_ddl\_time  
from all\_objects  
where owner = 'OUTLN'  
and object\_type IN ('TABLE', 'INDEX')  
**group by** owner, object\_type,  
to\_char(last\_ddl\_time, 'yyyymmdd hh24miss')  
**order by** object\_type, last\_ddl\_time;  
ObjTyp LAST\_DDL\_TIME

```
-----  
INDEX 20031001 173156  
INDEX 20080906 102159  
TABLE 20080906 102610
```

- **2 of 3 have same LAST\_DDL\_TIME at detail level, we'll use as demo for RANK**



# ROW\_NUMBER and RANK

```
select object_type "ObjTyp", substr(object_name,1,10) "ObjName",  
       to_char(last_ddl_time,'yyyymmdd hh24miss') last_ddl_time,  
       row_number() over  
         (partition by object_type order by last_ddl_time) RN,  
       rank() over  
         (partition by object_type order by last_ddl_time) R,  
       dense_rank() over  
         (partition by object_type order by last_ddl_time) DR  
from all_objects  
where owner = 'OUTLN' and object_type IN ('TABLE','INDEX');
```

Note that all three functions have same PARTITION BY and ORDER BY clauses.

Results on next slide ...



# Row\_number, Rank, Dense Rank

- Remember: Last three columns all have:  
(partition by object\_type order by last\_ddl\_time)  
[...]

ObjTyp	ObjName	LAST_DDL_TIME	RN	R	DR
INDEX	OL\$NAME	20031001 173156	1	1	1
INDEX	OL\$HNT_NUM	20031001 173156	2	1	1
INDEX	OL\$SIGNATU	20031001 173156	3	1	1
INDEX	OL\$NODE_OL	20080906 102159	4	<b>4</b>	<b>2</b>
TABLE	OL\$NODES	20080906 102610	1	1	1
TABLE	OL\$	20080906 102610	2	1	1
TABLE	OL\$HINTS	20080906 102610	3	1	1



# Behavior: Row\_number, Rank, Dense Rank

- If two records have the same value in the ORDER BY, the two records get different ROW\_NUMBER. RANK and DENSE\_RANK do not work like that.
- If two records have the same value in the ORDER BY, they both get the same RANK or DENSE\_RANK. The difference between RANK and DENSE\_RANK is how they are counted. DENSE\_RANK uses sequential numbers, RANK does not. In the INDEX set, the fourth is different from the first three. RANK jumps to display "4", and DENSE RANK is "2".



# ROW\_NUMBER is similar to ROWNUM

- One key difference:  
ROWNUM gets incremented as rows are returned from the query,  
so we can not say "WHERE ROWNUM = 5".  
But ROW\_NUMBER **can** be used that way.



# Sort both ways in same SQL

```
➤ Select [...],  
    row_number() OVER ( partition by object_type  
                        order by last_ddl_time) SORTUP,  
    row_number() OVER ( partition by object_type  
                        order by last_ddl_time DESC NULLS LAST) SORTDOWN  
from all_objects where owner = 'OUTLN'  
and object_type IN ('TABLE', 'INDEX');
```

ObjTyp	LAST_DDL_TIME	<b>SORTUP</b>	<b>SORTDOWN</b>
INDEX	20031001 173156	1	4
INDEX	20031001 173156	2	3
INDEX	20031001 173156	3	2
INDEX	20080906 102159	4	1
TABLE	20080906 102610	1	3
TABLE	20080906 102610	2	2
TABLE	20080906 102610	3	1



# Traditional -- slow way to see mixed detail and summary levels (generated from reporting tool)

```
> SELECT <detail columns>, max_effdt, max_effseq
FROM
  ( SELECT <detail columns>,
    MAX (DISTINCT t2.APLAN_EFFDT) max_effdt
  FROM t3, t1 LEFT OUTER JOIN t2 ON [...]
  WHERE [...] GROUP BY t1.CTERM_EMPLID, t1.CTERM_TERM_CD) d5,
  ( SELECT <detail columns>,
    MAX (t2.APLAN_EFFSEQ) max_effseq
  FROM t3, t1 LEFT OUTER JOIN t2 ON [...]
  WHERE [...] GROUP BY t1.CTERM_EMPLID, t1.CTERM_TERM_CD,
    t2.APLAN_EFFDT) d4,
  ( SELECT <detail columns only, no aggregate!!!>
  FROM t3, t1 LEFT OUTER JOIN t2 ON [...]
  WHERE [...] < NO group by clause!!!> ) d3
WHERE < predicates for outer select >
ORDER BY < columns for outer select >
```



# Analytic Function is FASTER!!



Requires only one inline view, a single pass instead of three...





# Improvement is Hours to Minutes

```
➤ SELECT <detail columns>, max_effdt, max_effseq
FROM
  ( SELECT <detail columns>,
    max(t2.APLAN_EFFDT) OVER
      (PARTITION BY t1.ctrm_emplid, t1.ctrm_term_cd)
      AS max_effective_date,
    max(t2.APLAN_EFFSEQ) OVER
      (PARTITION BY t1.ctrm_emplid, t1.ctrm_term_cd,
        t2.APLAN_EFFDT)
      AS max_effective_sequence
    FROM t3, t5 (t1 LEFT OUTER JOIN t2 ON [...] )
      LEFT OUTER JOIN t4 ON [...] WHERE [...] )
WHERE < predicates for outer select >
ORDER BY < columns for outer select >
```



# Keep adapting, don't be a dinosaur!



# Example of UPDATE

```
➤ UPDATE FZBRFCX SET FZBRFCX_ZERO_FLAG = 2
WHERE rowid IN
  (SELECT rowid FROM
    (SELECT rowid, FZBRFCX_ZERO_FLAG Flag,
      sum(FZBRFCX_TRANS_AMT) OVER
        (PARTITION BY FZBRFCX_ACCT_CODE,
          FZBRFCX_FUND_CODE,
          FZBRFCX_DOC_REF_NUM) Sum_Amt
    FROM FZBRFCX )
  WHERE Sum_Amt = 0 AND Flag <> 0 ) ;
```



# Helpful sidetrack: Query Subfactoring AKA **Common Table Expression**

- Analytic functions often need an inline view (subquery).
- Sometimes the inline views are nested
- Indentation is helpful, but can be confusing
- Subfactoring used here for clarity with related SQL statements across multiple slides
- Subfactoring easily allows multiple use of alias [11.2 allows recursive too]



# Traditional inline: layers with indentation

```
> SELECT MID_LVL.po_code, MID_LVL.seq, [...]
FROM
  (select INNER_LVL.po_code, INNER_LVL.seq, [...]
   from
     (select po_code, seq, [...]
      from fprpoda b
      where po_code in
        (select b.po_code
         from fprpoda b
         where activity_date
           between '01-NOV-09' and '09-NOV-09') CODE_LIST
     ) INNER_LVL
   ) MID_LVL
WHERE < [MID_LVL.column] predicates...>
```



# Query Subfactoring: Top Down

➤ WITH

```
CODE_LIST AS
```

```
(select po_code  
  from fprpoda  
  where activity_date  
         between '01-NOV-09' and '09-NOV-09' ),
```

```
INNER_LVL AS
```

```
(select po_code, seq, [...]  
  from fprpoda  
  where po_code in CODE_LIST ),
```

```
MID_LVL AS
```

```
(select po_code, seq, [...]  
  from INNER_LVL )
```

```
SELECT * FROM MID_LVL  
WHERE < predicates...>
```



# Running Totals and Windowing

- Requirement:  
Show values from current and previous rows, where running total went above \$50,000 level, where more stringent approvals are required:  
Is anyone trying to get around audit rules?
- Originally looked like it would need PL/SQL, with cursors starting and stopping
- We'll use query subfactoring to see the pieces build on each other...



# Running Totals and Windowing: stmt1

## ➤ WITH

```
code_list AS ( -- [codes used in next stmt]
SELECT distinct po_code
FROM fprpoda
WHERE trunc(activity_date)
      BETWEEN '01-NOV-09' AND '09-NOV-09'
      AND seq is not null ),
```





# Running Totals and Windowing: stmt2

```
➤ INNER_LVL AS ( -- [sum each code and seq combo]
SELECT po_code,
       LAG(po_code, 1) OVER
         (ORDER BY po_code) "PrevCode",
       seq , amt "CurrAmt", activity_date,
       SUM(amt) OVER (PARTITION BY po_code
                     ORDER BY po_code, seq, activity_date)
         running_tot
FROM fprpoda
WHERE po_code IN
( select po_code from CODE_LIST ) ),
```



# Running Totals and Windowing: stmt3

```
➤ MID_LVL AS ( -- get curr/prev row values
SELECT      po_code, seq ,
            (CASE WHEN "PrevCode" != po_code THEN NULL
             ELSE  LAG(running_tot, 1) OVER
                    (ORDER BY po_code, seq) END)  "PrevRunTot",
            running_tot "RunningTot",
            activity_date curr_actv,
            (CASE WHEN "PrevCode" != po_code THEN NULL
             ELSE  LAG(activity_date) OVER
                    (ORDER BY po_code, seq) END)  prev_actv
FROM INNER_LVL)
-- 1 is default for LAG, hard coding would be
-- for clarity
```



# Running Totals and Windowing: Final

- Query subfactoring above is done: one isolated stmt shows what we're ultimately trying to do...

```
SELECT po_code, seq, "PrevRunTot",  
       "RunningTot"-"PrevRunTot" "DiffChange",  
       "RunningTot" , prev_actv, curr_actv  
FROM MID_LVL  
WHERE "PrevRunTot" < 50000  
      AND "RunningTot" >= 50000;
```

PO_CODE	SEQ	PrevRunTot	DiffChange	RunningTot	PREV_ACTV	CURR_ACTV
<b>B0142584</b>	<b>7</b>	<b>46,800.00</b>	<b>5,500.00</b>	<b>52,300.00</b>	<b>05-FEB-09</b>	<b>05-NOV-09</b>
<b>B0181676</b>	<b>1</b>	<b>38,142.00</b>	<b>23,856.34</b>	<b>61,998.34</b>	<b>26-NOV-07</b>	<b>17-NOV-08</b>
<b>S0176940</b>	<b>1</b>	<b>43,371.00</b>	<b>42,156.00</b>	<b>85,527.00</b>	<b>17-JUN-05</b>	<b>23-MAR-06</b>
<b>S0181330</b>	<b>1</b>	<b>1.00</b>	<b>302,069.91</b>	<b>302,070.91</b>	<b>20-JUL-07</b>	<b>28-AUG-07</b>



Detail for one PO, Seq# 0, 2, and 7: exact same date/time,  
so running total is not gradually increasing

PO_CODE	SEQ	CurrAmt	RunningTot	Activity_Date_Time
B0142584	0	1.00	5,001.00	22-JAN-2003 10:27:00
B0142584	0	5,000.00	5,001.00	22-JAN-2003 10:27:00
B0142584	1	6,500.00	11,501.00	09-OCT-2003 14:36:01
B0142584	2	-1.00	18,500.00	27-OCT-2004 15:51:01
B0142584	2	7,000.00	18,500.00	27-OCT-2004 15:51:01
B0142584	3	9,500.00	28,000.00	05-OCT-2006 13:27:01
B0142584	4	4,000.00	32,000.00	25-OCT-2007 09:45:02
B0142584	5	5,500.00	37,500.00	27-NOV-2007 10:12:03
B0142584	6	9,300.00	46,800.00	05-FEB-2009 11:12:01
B0142584	7	-7,000.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	7,000.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	-9,300.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	9,500.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	-4,000.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	-11,500.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	9,500.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	-9,500.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	11,500.00	52,300.00	05-NOV-2009 12:27:01
B0142584	7	9,300.00	52,300.00	05-NOV-2009 12:27:01



# Running Totals and Windowing: Notes

- "LAG" puts curr/prev values on same row, that allows easy WHERE clause to find threshold
- We could not put "LAG" in stmt with running total (needed extra layering), because curr/prev row was not available until running total was done
- We got new running total for each code, because that is in "PARTITION BY" clause



# 11.2 feature: LISTAGG

- Can be Simple Aggregate OR Analytic
- Concatenates values from rows into a string, i.e. a LIST AGGregation
- Example is continuation of Running Total, which has duplicate dates for some Sequences, and LISTAGG faithfully shows all dups
- Including “distinct” in SQL looks *across* column values, not *within* LISTAGG: can not eliminate dups
- Simple Aggregate example shows dups...



# 11.2 LISTAGG: Simple Aggregate

```
SELECT seq,  
       LISTAGG(to_char(activity_date, 'MON-YYYY'), '; ' )  
       WITHIN GROUP (ORDER BY seq) "Activity_Dates"  
FROM fprpoda WHERE po_code = 'B0142584' AND seq < 3  
GROUP BY seq;  
SEQ Activity_Dates
```

```
-----  
0 JAN-2003; JAN-2003  
1 OCT-2003  
2 OCT-2004; OCT-2004
```



## 11.2 LISTAGG: Analytic

- Need "distinct" for analytic, else shows all 19 rows;  
No GROUP BY, analytic function is at detail level

```
SELECT distinct seq,  
SUM(amt) OVER (ORDER BY seq, activity_date) "RunTot",  
SUM(amt) OVER (PARTITION BY seq ORDER BY  
seq, activity_date) "SeqTot",  
LISTAGG(amt, ';' ) WITHIN GROUP (ORDER BY seq)  
OVER (PARTITION BY seq) "Amts"  
FROM fprpoda WHERE po_code = 'B0142584'  
AND seq IS NOT NULL ORDER BY seq;
```

- Sequence 7 has 10 Amount entries, some cancel each other out...





# 11.2 LISTAGG: Analytic

➤ Reminder of syntax, and result from SQL stmt:

```
LISTAGG(amt, ';' ) WITHIN GROUP (ORDER BY seq)  
OVER (PARTITION BY seq) "Amts"
```

```
Seq RunTot SeqTot Amt
```

Seq	RunTot	SeqTot	Amts
0	5001	5001	5000;1
1	11501	6500	6500
2	18500	6999	7000;-1
3	28000	9500	9500
4	32000	4000	4000
5	37500	5500	5500
6	46800	9300	9300
7	52300	5500	9300;-9500;9500;-11500;-4000; 9500;-9300;7000;-7000;11500



# Moving Average

```
➤ WITH SGLCODE AS (  
  select seq, amt  
  from fprpoda  
  where po_code in ('BA177629')  
        and seq IS NOT NULL )  
SELECT   seq, amt,  
        avg(amt) OVER (order by seq rows  
                      between 1 preceding and 1 following ) ma1,  
        avg(amt) OVER (order by seq rows  
                      between 0 preceding and 1 following ) ma2,  
        avg(amt) OVER (order by seq rows  
                      between 1 preceding and 0 following ) ma3  
FROM SGLCODE order by seq;
```



# Moving Average: Result

➤ SEQ	AMT	MA1	MA2	MA3
0	.01	1068.86	1068.86	.01
0	2137.70	2379.24	3568.85	1068.86
1	5000.00	2596.23	2825.50	3568.85
2	651.00	18500.00	25250.00	2825.50
3	49849.00	25250.00	49849.00	25250.00

## NOTE:

MA1: 1 before, 1 after (3 rows avg)

MA2: 0 before, 1 after (2 rows avg)

MA3: 1 before, 0 after (2 rows avg)



# (Finally) NTILE

- Example on next slide is 6 rows of test scores with 4 buckets (Quartile)
- NTILE definition is ordered DESCENDING so that highest test scores are in buckets 1 and 2
- If “DESC” were taken out of SQL, ranking would be reversed, i.e. lowest scores would be in the 1st quartile rather than 4th
- Two extra values (  $6/4$  ) are allocated to buckets 1 and 2



# NTILE: Example with 4 buckets

```
SELECT name, score,  
       NTILE(4) OVER (ORDER BY score DESC)  
       AS quartile  
FROM test_scores ORDER BY name;
```

NAME	SCORE	QUARTILE
Barry Bottomly	12	4
Felicity Fabulous	99	1
Felix Fair	41	2
Mildred Middlin	55	2
Paul Poor	24	3
Sharon Swell	86	1



# A & Q



- **A & Q**      **Answers: Wisdom to share?**      **Questions?**
- Philip Rice    price [at] ucsc {dot} edu

